

Barabási Queueing Model and Invasion Percolation on a tree

A. Gabrielli¹, G. Caldarelli¹

¹SMC, INFN-CNR, Dipartimento di Fisica Università di Roma "Sapienza", Piazzale A. Moro 2, 00185 Rome, Italy, ISC, CNR, Via dei Taurini 19, 00185 Rome, Italy.

In this paper we study the properties of the Barabási model of queueing under the hypothesis that the number of tasks is steadily growing in time. We map this model exactly onto an Invasion Percolation dynamics on a Cayley tree. This allows to recover the correct waiting time distribution $P_W(\tau) \sim \tau^{-3/2}$ at the stationary state (as observed in different realistic data) and also to characterize it as a sequence of causally and geometrically connected bursts of activity. We also find that the approach to stationarity is very slow.

PACS numbers: 89.75.Da, 02.50.Le, 89.65.Ef

Queueing theory [1, 2, 3] describes a wide range of human dynamical behaviors [4, 5]. Most of traditional models lead to exponential *waiting time distributions* (WTD) $P_W(\tau) \sim \exp(-\tau/\tau_0)$ for the tasks in the queue. Recently, motivated by observations related to web browsing, email communications and ordinary mail correspondence [6], much attention has been paid to priority driven queueing models generating power-law WTD $P_W(\tau) \sim \tau^{-\alpha}$ for the tasks. In this paper we study a particular version of one of the latter kind: the Barabási queueing model (BQM) [7, 8]. In our version of BQM at each time step the task with highest random priority is always executed and replaced in the queue by a constant number $m \geq 2$ of new tasks with random priorities. This process can be mapped exactly onto an Invasion Percolation (IP) dynamics [9] on a Cayley tree [10] with a series of advantages. Firstly we can characterize the task list dynamics through the WTD at the stationary state. Secondly we show that its general evolution is composed by a sequence of geometrically and causally connected burst of activities (task avalanches) with scale-invariant size distribution. Thirdly, we can study the dynamics out of stationarity and we show that the approach to it is very slow. Finally, it permits to simply generalize the results in to the case of time-varying m . In the general BQM [7] one starts with an initial list (i.e. queue) of $n_0 \geq 2$ tasks. At every time-step t one of these tasks is executed and replaced by $m(t)$ other new tasks. For constant $m(t) = 1$, the queue length remains constant. The execution rule at each time-step is given by fixing a random priority index $x_i \in [0, 1]$ for each task in the queue and then executing with a probability $p \leq 1$ the task with the highest priority and with a probability $(1-p)$ a randomly chosen task. The related problem for general $0 \leq p \leq 1$ and $m = 1$ has been analyzed and solved in [11, 12]. In the purely extremal (i.e. when $p = 1$) case with a variable queue length, the behavior of $P_W(\tau)$ differs strongly from the previous case. In [13] it has been studied the case in which at each time-step there is a probability $\mu \leq 1$ to execute the highest priority task, while a new task is added to the list with another probability $\nu \leq 1$. For $\mu = \nu = 1$ the above case of conserved queue length is recovered. If at least one among μ and ν is strictly smaller than 1, the list

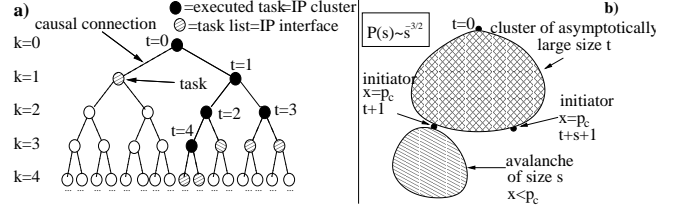


FIG. 1: a) Sketch of the first four steps of IP dynamics on a Cayley tree with branching ratio $m = 2$. b) Illustration of a causally and geometrically connected avalanche whose size distribution is $P(s) \sim s^{-3/2}$. This characterizes the stationary state of both IP and the task queue dynamics.

length instead varies in time. Depending whether $\mu > \nu$, $\mu = \nu < 1$ or $\mu < \nu$ the WTD $P_W(\tau)$ at the stationary state changes the asymptotic behavior. In particular for $\mu = \nu < 1$ all tasks are executed with $P_W(\tau) \sim \tau^{-3/2}$ with no upper cut-off, while for $\mu < \nu$, the mean queue length grows linearly in time, and one can show that asymptotically for $t \rightarrow \infty$ all tasks with priority index $x_i < (1 - \mu/\nu)$ are never executed staying forever in the queue. Instead tasks with $x_i \geq (1 - \mu/\nu)$ are executed with a WTD coinciding with the one for $\mu = \nu < 1$. In our version of the model the most urgent task is executed with probability $\mu = 1$ and replaced in the list by a constant number $m \geq 2$ of new tasks with random priorities. All the features of this model can be clarified by mapping it into an IP dynamics on a Cayley tree. A similar mapping on 1-d IP has proved to be fruitful also in the case of fixed queue length [12]. Invasion Percolation on a Cayley tree [10] is defined as follows (see Fig. 1-a): let us take a Cayley tree with branching ratio m where initially only the top vertex site of the tree is occupied. A random number (*fitness*) x , extracted from a given probability density function (PDF) $p(x)$, is assigned once and forever to each empty site (independently of the others). At each time-step the site of the growth interface ∂C_t with the highest fitness is occupied. ∂C_t is defined at each time t as the set of empty sites connected by a first nearest neighbor rule to the connected growing cluster C_t of occupied sites up to that time. Since for each occupied site other new m sites enter the growth interface, the number of sites

respectively in \mathcal{C}_t and in $\partial\mathcal{C}_t$ at time t are respectively $\|\mathcal{C}_t\| = (t+1)$ and $\|\partial\mathcal{C}_t\| = m + (m-1)t$. Being the dynamics extremal, the statistical and geometrical features of IP are independent on the shape of $p(x)$; our choice is to take $p(x) = 1$ with $x \in [0, 1]$.

The exact mapping between IP and our queueing model is done by identifying sites with tasks, fitness with priority index, growth interface $\partial\mathcal{C}_t$ in IP with the task list (i.e. the queue), and finally the growing IP cluster \mathcal{C}_t with the set of executed tasks up to time t .

For our purposes we focus on the following features of the asymptotic stationary state of IP dynamics:

1) The distribution (also called *normalized interface histogram*) $\phi_s(x)$ of the fitnesses of the interface sites (i.e. of the tasks in the queue), has the step-function shape

$$\phi_s(x) = p_c^{-1} \theta(p_c - x),$$

where $p_c = (m-1)/m$ is the ordinary percolation threshold of the Cayley tree. This implies that: (i) apart from a vanishing fraction (i.e. a finite number) of sites, all the interface sites have fitness $x < p_c$; (ii) since the number of sites in the stationary state is infinite, only those few sites with $x \geq p_c$ can grow at each time-step. Indeed, at each time for the just occupied site (executed task) $m \geq 2$ new sites (new tasks) enter the interface (queue). This implies that a fraction $(m-1)/m$ of the sites entered the interface at any time will never be executed. Since the interface site with maximal fitness is always executed and the “fresh” interface sites have random fitness, asymptotically only and all the sites with $x \geq (m-1)/m$ are executed while the others stay forever in the interface;

2) The cluster of occupied sites is substantially coinciding with the incipient percolating cluster of ordinary percolation (i.e. at occupation probability $p = p_c$);

3) The stationary dynamics self-organizes into a sequence of spatially and causally connected avalanches of growth activity [14] with a scale-invariant size distribution $P(s) \sim s^{-3/2}$ independently of the value of m . Any of these avalanches (see Fig. 1-b), say \mathcal{A} , starts with the growth of a site (the *initiator* of \mathcal{A}) with fitness $x = p_c$ exactly, meaning that all the other interface sites at that time have $x \leq p_c$. Following this growth, m new sites/tasks (*sons*), geometrically connected to the initiator, enter the interface. \mathcal{A} stops immediately if all sons have fitness $x < p_c$, and consequently another “old” interface site grows with $x = p_c$ [due to the shape of $\phi_s(x)$], and therefore initiating a new avalanche \mathcal{B} , i.e., \mathcal{A} lasted only one step. If instead at least one of the sons of the initiator of \mathcal{A} has $x > p_c$, then \mathcal{A} goes on at least one step further as one of these sons grows. Consequently, other new m “descendants” (sons of a son) of the initiator of \mathcal{A} enter the interface. Again \mathcal{A} keeps on if at least one among all the remaining descendants of any generation (called the *avalanche interface*) has fitness $x > p_c$ otherwise the avalanche stops, and so on.

The exponent of $P(s)$ can be computed analytically by mapping the avalanche dynamics into a problem of first return of unbiased random walks. Let n_t be the number

of sites (tasks) with $x \geq p_c = (m-1)/m$ after the t^{th} step of an avalanche (i.e., on its interface). Since after the growth of one site m new sites enter the interface, we have the following Markovian evolution for $n_t > 0$

$$n_{t+1} = n_t + j - 1 \quad \text{with prob.} \quad \binom{m}{j} p_c^{m-j} (1-p_c)^j \quad (1)$$

with $j = 0, 1, \dots, m$. I.e., n_t follows an ordinary random walk with independent steps. As $p_c = (m-1)/m$, the average increment of n_t in one time-step is zero (*martingale property*). Therefore [15] the probability distribution of the time s for which $n_s = 0$ for the first time (i.e. the duration of the avalanche) scales as $s^{-3/2}$ at large s . A percolation argument can also be used to find out the same exponent: as the avalanche initiator has $x = p_c$ and the avalanche lasts exactly for a time interval equal to the number of sites with $x > p_c$ connected to it in the positive time direction, then the avalanche size is distributed as the finite clusters at the critical point $p = p_c$ in ordinary percolation on the same tree: $P(s) \sim s^{-3/2}$.

Random walk and diffusion theory arguments also permit to evaluate the stationary state WTD $P_W(\tau)$ for the tasks with $x > p_c$. We follow here a similar discussion to [13]. We can write the WTD as

$$P_w(\tau) = \sum_{n=0}^{\infty} \int_{p_c}^1 dx \tilde{Q}(n, x) G(n, x, \tau) \quad (2)$$

where $\tilde{Q}(n, x)$ is the probability that at a generic time-step at the stationary state we have exactly n tasks in the queue (i.e. sites on the IP interface) with priority larger than $x \geq p_c$. The quantity $G(n, x, \tau)$ is instead the conditional probability that, always at the stationary state, a certain task with priority $x \geq p_c$ added to the list at a time-step when other n tasks with priority larger than x are present, is executed after τ time-steps. We can write the evolution equation for the the number $n_t(x)$ of tasks in the list with priority larger than x at time t . Similarly to Eq. (1) we can write for $n_t(x) \geq 1$

$$n_{t+1}(x) = n_t(x) + j - 1 \quad \text{with prob.} \quad \binom{m}{j} x^{m-j} (1-x)^j, \quad (3)$$

where $j = 0, 1, \dots, m$. We can consequently write the master equation for the probability $Q(n, x, t)$ that at time t there are exactly n tasks in the list with priority larger than x . To aim of simplicity let us write it for $m = 2$ for which $p_c = 1/2$. From Eq. (3) we can write for $n \geq 3$

$$Q(n, x, t+1) = Q(n+1, x, t)x^2 + Q(n, x, t)2x(1-x) + Q(n-1, x, t)(1-x)^2 \quad (4)$$

while for $n \leq 2$ we have

$$\begin{aligned} Q(2, x, t+1) &= Q(3, x, t)x^2 + Q(2, x, t)2x(1-x) \\ &\quad + Q(1, x, t)(1-x)^2 + Q(0, x, t)(1-x)^2 \\ Q(1, x, t+1) &= Q(2, x, t)x^2 + Q(1, x, t)2x(1-x) \\ &\quad + Q(0, x, t)2x(1-x); \\ Q(0, x, t+1) &= Q(1, x, t)x^2 + Q(0, x, t)x^2 \end{aligned} \quad (5)$$

$\tilde{Q}(n, x)$ is given by the stationary solution of the above equations. In order to find both $\tilde{Q}(n, x)$ and $G(n, x, \tau)$ we can now proceed in a similar way to [13]. It is simple to show that the well-normalized stationary solution for $x \geq p_c = 1/2$ of Eqs. (4) and (5) is

$$\begin{aligned}\tilde{Q}(n, x) &= \frac{2(x - p_c)}{x^2} \left[\frac{(1 - x)^2}{x^2} \right]^{n-1} \quad \text{for } n \geq 2 \quad (6) \\ \tilde{Q}(1, x) &= 2 \frac{1 - x^2}{x^2} (x - p_c); \quad \tilde{Q}(0, x) = 2(x - p_c).\end{aligned}$$

Note that for $x \rightarrow p_c^-$ any $\tilde{Q}(n, x) \rightarrow 0$ with the ratio $\tilde{Q}(n, x)/\tilde{Q}(l, x) \rightarrow 1$ for any $n, l \geq 2$, i.e., the distribution of the number $n_{t \rightarrow \infty}(p_c)$ becomes practically uniform.

The quantity $G(n, x, \tau)$ can be found by Eq. (4) in complete analogy with [13] and [16] leading both to the same correct scaling behavior $P_W(\tau) \sim \tau^{-3/2}$. We refer here to the former as it is of simpler formulation. First of all we note that Eq. (4), in both the continuous time and $n = y$ approximation, becomes the diffusion equation:

$$\partial_t Q(y, x, t) = c(x) \partial_y^2 Q(y, x, t) + d(x) \partial_y Q(y, x, t), \quad (7)$$

with $c(x) = x^2$ and $d(x) = x^2 - (1 - x)^2$. Since we are considering $x \geq p_c = 1/2$ we have $d(x) \geq 0$, i.e., there is a drift to the small y (i.e. n) direction. $G(n, x, \tau)$ can be seen as the probability that at the stationary state, fixed x and given that at time $t = 0$ it is $y = n$, one has $y = 0$ for the first at time $t = \tau$. This implies that [13, 15]

$$G(n, x, t) = -\partial_t \left[\int_0^\infty dy Q(y, x, t) \right],$$

where here $Q(y, x, t)$ is the solution of Eq. (7) with initial condition $Q(y, x, 0) = \delta(y - n)$. All this gives

$$G(n, x, \tau) = \frac{n}{\sqrt{4\pi c(x)t}} \exp \left\{ -\frac{[n - d(x)t]^2}{4c(x)t} \right\}$$

We now use this result and Eq. (6) in Eq. (2) to find $P_W(\tau)$. It is simple [13] to show that for large τ we have $P_W(\tau) \sim \tau^{-3/2}$. In other words each task with $x \geq p_c$ has to wait a finite portion of the avalanche duration before being executed. Note that all these results are completely independent of the integer branching factor $m > 1$. From Eq. (3) it is natural to expect to have the same result in the case in which at each time step m is not constant but fluctuates with independent fluctuations such that $\langle m \rangle \geq 1$ and finite variance. This is the reason why our model share the same statistical features with that in [13]. In the case where $\langle m^2 \rangle = +\infty$ we expect anomalous exponents for both $P(s)$ and $P_W(\tau)$ as the random walks Eqs. (1) and (3) become Levy flights as shown in [17].

We now address the question on how fast is the approach to stationarity in such models. Again some rigorous theoretical results for IP on a tree turn to be very useful to this end. We summarize here the main results in literature, and then propose a simple mean-field approach showing how slow the relaxation to the right stationary state is. In [10] the main exact result, adapted to

our notation, states that the probability that at time t of the dynamics a task with priority smaller than $(p_c - \epsilon)$ is executed, vanishes exponentially fast for large t for $\epsilon > 0$, but as $t^{-1/2}$ for $\epsilon \rightarrow 0^+$. This suggests that deviations from the stationary dynamics disappear as $t^{-1/2}$. In [18] instead it is shown rigorously that: (i) IP on a Cayley tree has in the infinite time limit a unique backbone. In terms of the task dynamics this means that there is a unique infinite chain of executed task which are causally connected in the IP sense above. (ii) The minimal priority of the executed tasks staying on the backbone beyond the k^{th} generation of the Cayley tree (see Fig. 1-a) is $p_c(1 - Z/k)$ for large k where Z is an exponential random variable with unitary mean.

We now present a simple mean field argument showing this slow approach of the list dynamics to the right stationary state. We study the dynamics of the above introduced normalized distribution $\phi(x, t)$ of the priorities of the tasks in the queue (fitness *histogram* of interface sites in IP) at time t . In order to write a closed equation for $\phi(x, t)$ we use the Run Time Statistics (RTS) which is a probabilistic method introduced to describe IP and related dynamics, and evaluate the statistical weight of all different “histories” of the dynamics (i.e. paths in the realization space) [19, 20, 21]. Let $h(x, t)dx$ be the number of tasks in the queue at time t with priority in $[x, x + dx]$ in a single realization. We can write rigorously:

$$h(x, t + 1) = h(x, t) - m(x, t + 1) + m, \quad (8)$$

where $m(x, t)$ is the PDF of the priority of the executed task at time t conditional to the whole past history. By calling $p_i(x, t)$ the PDF of the priority of the i^{th} task in the queue at time t conditional to the same past history, and assuming that the executed task at that time is the j^{th} , a good approximation for $m(x, t + 1)$ [20] is $m(x, t + 1) = \frac{1}{\mu_j(t)} p_j(x, t) \prod_{i \neq j}^{\partial \mathcal{C}_t} \left[\int_0^x dy p_i(y, t) \right]$, where $\mu_j(t) = \int_0^1 dx p_j(x, t) \prod_{i \neq j}^{\partial \mathcal{C}_t} \left[\int_0^x dy p_i(y, t) \right]$ is the probability of selecting j conditional to the past history. We now average Eq. (8) over all the possible realizations of the dynamics up to time $(t + 1)$ using the symbol $\langle \cdot \rangle_{t+1}$ for this average. By definition we have $\langle h(x, t + 1) \rangle_{t+1} = (m - 1)(t + 1)\phi(x, t + 1)$ and $\langle h(x, t) \rangle_{t+1} = \langle h(x, t) \rangle_t = (m - 1)t\phi(x, t)$. In order to take the same average of $m(x, t + 1)$ note that, if $A(i_0, i_1, \dots, i_{t-1}, i_t)$ is a function of the queue history up to time $(t + 1)$ identified by the sequence of executed tasks $\{i_0, i_1, \dots, i_{t-1}, i_t\}$, by the rules of conditional probability, we can write

$$\langle A(i_0, i_1, \dots, i_t) \rangle_{t+1} = \left\langle \sum_{j \in \partial \mathcal{C}_t} \mu_j(t) A(i_0, i_1, \dots, i_{t-1}, j) \right\rangle_t.$$

We therefore have

$$\langle m(x, t + 1) \rangle_{t+1} = \left\langle \sum_{j \in \partial \mathcal{C}_t} p_j(x, t) \prod_{i \neq j}^{\partial \mathcal{C}_t} \left[\int_0^x dy p_i(y, t) \right] \right\rangle_t.$$

Considering that by definition $\langle p_j(x, t) \rangle_t = \phi(x, t)$, we now introduce the mean field approximation consisting

in replacing the average of the above products of $p_i(x, t)$ with the products of the averages, i.e.,

$$\langle m(x, t+1) \rangle_{t+1} = (m-1)t\phi(x, t) \left[\int_0^x dy \phi(y, t) \right]^{(m-1)t-1}. \quad (9)$$

We can now write the mean-field equation for $\phi(x, t)$ as

$$\begin{aligned} \phi(x, t+1) &= \frac{t}{t+1} \phi(x, t) \left\{ 1 - \left[\int_0^x dy \phi(y, t) \right]^{(m-1)t-1} \right\} \\ &+ \frac{m}{(m-1)(t+1)} \end{aligned} \quad (10)$$

This strong decorrelating approximation is expected to lead to a faster relaxation to stationarity than the actual one. We show however that, even in this approximation, the stationary state is the right one and the approach to it is power law. Integrating both sides of Eq. (10), taking the continuous time limit and $t \gg 1$ we get

$$\partial_t \psi(x, t) = \frac{-1}{t+1} \left[\psi + \frac{1}{m-1} \psi^{(m-1)t} - \frac{mx}{m-1} \right] \quad (11)$$

where $\psi(x, t) = \int_0^x dx' \phi(x', t)$ is the cumulative average priority distribution, and we have assumed $\psi(0, t) = 0$ at all t . The initial condition for Eq. (11) is $\psi(x, 0) = x$. Since $\phi(x, t)$ is a normalized PDF, we have $\psi(x, t) \geq 0$, non-decreasing in x and $\psi(1, t) = 1$.

In the x region were $(1 - \psi) \gg 1/[(m-1)t]$ we can approximate Eq. (11) simply with

$$\partial_t \psi(x, t) = -\frac{1}{t+1} \left(\psi - \frac{mx}{m-1} \right), \quad (12)$$

which leads to the solution for sufficiently large t

$$\psi(x, t) = \frac{mx}{m-1} \left(1 - \frac{1}{mt} \right) \text{ for } x < \frac{m-1}{m} - \frac{1}{m^2 t}. \quad (13)$$

Note that $p_c = (m-1)/m$. Moreover in the x region were $\epsilon = (1 - \psi) \ll 1/[(m-1)t]$ it is simple to show that the following approximation holds

$$\partial_t \epsilon(x, t) = -\epsilon(x, t) + \frac{m(1-x)}{(m-1)(t+1)} \quad (14)$$

whose solution is $\psi(x, t) = 1 - \epsilon(x, t)$ with

$$\epsilon(x, t) = \frac{m(1-x)}{(m-1)(t+1)} \left[1 + O\left(\frac{1}{t}\right) \right] \quad (15)$$

when $x \gg (m-1)/m$. All this means that

$$\phi(x, t) = \frac{\theta(p_c - x)}{p_c} + \delta\phi(x, t)$$

with $\delta\phi(x, t) \sim 1/t$. Therefore even in this mean field approximation, for which we expect a faster relaxation, deviations from it vanish as slowly as $1/t$.

In conclusion, we have shown a way to analytically compute all the main features of the Barabási model of human dynamics with time-increasing queue length. This is done by using Invasion Percolation on a Cayley tree and random walk theory. We believe that the approach we introduced, allows us to describe quantitatively two intuitive features of tasks queues. The first feature is that some tasks seem to remain indefinitely before being processed; secondly we recover naturally the fact that in real world execution of a task has often the effect to generate an avalanche of new tasks. Through our approach one can study both the stationary state dynamics and the approach to it. This shows that both are characterized by temporal power laws as typical for extremal dynamics in quenched disorder [19, 20, 21].

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